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Methods and Systems For Protecting Media Content

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TECHNICAL FIELD

This invention relates to methods and systems for protecting media content.

BACKGROUND

Content providers such as those that provide audio and/or video data in the form of data streams, application vendors that provide applications to render such data, and others often desire to protect such data from being ascertained, used or otherwise rendered by unauthorized entities. Such protection is typically desired from the point where the data resides, in some protected form and on some type of readable medium, along the chain of components that will handle or otherwise process such data, up to and including both the physical output connector that provides the data to a device such as a display or speakers, as well as these devices themselves.

As an example, consider Fig. 1 which illustrates, at a high level, an exemplary system 10 into which protected content 12, e.g. compressed and encrypted content, can be received and processed. In this example, protected content 12 is provided into a consumer device 14 such as a personal computer. Within the consumer device, an application 16 typically processes the protected content by, for example, using a decryption component 18 to decrypt the content and a decoder component 20 to decode or decompress the content. Once the content is decoded, it can typically undergo some processing, after which time the processed content is provided to a renderer component 22 which then ensures that the data is provided to a device, such as a display 24 (in the case of video data) and/or speakers 26 (in the case of audio data).

1 In this system, in order for the data to get to the appropriate hardware
2 device such as the display or speaker, the data has to transit some type of bus, such
3 as a PCI/AGP bus, and then has to travel through a physical connector and over a
4 cable, such as an S-Video cable or a DVI cable.

5 Over time, various different types of content protection schemes and
6 technologies have been developed and employed to protect the content when, for
7 example, the content is decrypted within the consumer device 14 and then
8 provided to an output device such as a display or speakers. These protection
9 schemes and technologies are as varied as the different kinds of audio and video
10 data. In addition, and perhaps what is most important is that today, there is no
11 way for an application to securely control the *behavior* of collections of these
12 various different types of content protection schemes and technologies.

13 More specifically, consider the chain of components that process data that
14 is to be ultimately rendered on a hardware device. As will be appreciated by the
15 skilled artisan, there are different layers of components that touch or otherwise
16 have access to this data and associated communications from the application while
17 the data and the communications are in the chain. For example, the application
18 typically does not communicate directly to the hardware that is to render the data.
19 Rather, there is typically a software stack in the user mode with which the
20 application communicates, and then a driver stack in the kernel mode below the
21 user mode software stack that actually communicates with the hardware.
22 Effectively, any one of these components in the chain can actually tamper with the
23 data and/or communications being passed along the chain.

1 Accordingly, this invention arose out of concerns associated with providing
2 secure channels for both communications and data to flow from the application to
3 downstream components.
4

5 **BRIEF DESCRIPTION OF THE DRAWINGS**

6 Fig. 1 is a high level block diagram of a system that can process media
7 content.

8 Fig. 2 is a block diagram of a system that is used to illustrate aspects of a
9 protection protocol in accordance with one embodiment.

10 Fig. 3 is a block diagram that illustrates aspects of the inventive protection
11 protocol in accordance with one embodiment.

12 Fig. 4 is a block diagram that illustrates aspects of the inventive protection
13 protocol in accordance with one embodiment.

14 Fig. 5 is a block diagram that illustrates aspects of the inventive protection
15 protocol in accordance with one embodiment.

16 Fig. 6 is a flow diagram that describes steps in a method in accordance with
17 one embodiment.

18 Fig. 7 is a block diagram that illustrates an exemplary system in connection
19 with an implementation example.
20

21 **DETAILED DESCRIPTION**

22 **Overview**

23 Various embodiments pertain to methods and systems that utilize a protocol
24 which enables media content protection by establishing a secure communication
25 channel and, in some embodiments, a secure data channel, between a device such

1 as a computer running a protected content playback application, and an associated
2 driver, such as a graphics driver, of an associated display device such as a monitor,
3 flat panel LCD, television and the like. Various embodiments are directed to
4 addressing the needs of content providers and application vendors to signal that
5 content protection should be applied to data that is output over a physical
6 connector to a device such as a display device.

7 The various embodiments described below provide a means for securely
8 passing commands from a user mode playback application to a driver, such as a
9 graphics driver, and securely returning status information from the driver to the
10 user mode application. In the embodiments described below, once a secure
11 communication channel is established, the user mode application can instruct an
12 associated driver to enable content output protection on the physical connector
13 between the computer and the display device. Exclusivity of the secure channel
14 session can prevent attackers from recording the protected content as it is
15 transmitted from the computer to the display.

16 In the discussion that follows, the description of the various embodiments is
17 directed along two separate but related lines. Specifically, in the section below
18 entitled "Protocol for Establishing a Secure Channel from the Application to
19 Downstream Components", a description of the inventive protocol is provided.
20 This protocol is intended to be used in connection with any suitable system and is
21 not necessarily tied to any one particular system or operating system. In the
22 section entitled "Exemplary Implementation of Protocol on a Windows®
23 Operating System", a description of the protocol is provided in the context of an
24 implementation that involves a Windows® operating system. This description is
25

1 intended to illustrate, among other things, one specific implementation and some
2 of the design considerations associated with this specific implementation.

3 Protocol for Establishing a Secure Channel from the Application to 4 Downstream Components

5
6 In the discussion that follows, at least some embodiments are described that
7 can be used to provide a secure signaling mechanism from an application to
8 various downstream components so that the application can communicate which
9 of a potential number of content protection technologies should be applied, and
10 how at least some of those content protection technologies should be applied to
11 protect media content output. In addition, and building on the notion of enabling
12 and configuring different types of content protection, at least some of the
13 embodiments are agnostic with respect to the *type* of content protection
14 technology that is used to protect the media output. This is advantageous in that
15 the described system and protocol are adaptively flexible and be used in
16 connection with many different types of current and future content protection
17 technologies, as will become apparent below.

18 As an example overview of a typical system in which the inventive protocol
19 can be used, consider Fig. 2 which shows a computing system in the form of a
20 consumer device 200. Consumer device 200 can comprise any suitable type of
21 consumer device that can be utilized to render protected media content. One
22 example of such a device is a personal computer. Other types of devices can
23 include personal video recorder (PVR) devices, digital video recorder (DVR)
24 devices and the like.
25

1 Consumer device 200 comprises a playback application 202 that can
2 acquire protected media content from a protected media content source, system
3 software 204 that facilitates processing of the protected media content, and driver
4 software 206 (or the software representation of associated hardware) that is
5 utilized to interface with hardware 208 (such as a video card) that provides the
6 processed media content over some type of physical connector 210 to an
7 associated output device 212 over which the media content is rendered for a user.
8 Examples of output devices include display devices and speakers, as noted above.

9 For purposes of the ongoing discussion, at least in this section, the *types* of
10 playback applications, system software, driver software and hardware are not
11 particularly relevant, as the skilled artisan will surely appreciate that there are
12 many various different types of entities that can make up the chain between a
13 playback application and the physical connector. As such, it is not the intent of
14 this discussion to be limited to any one particular type of system.

15 In the illustrated and described embodiment, there are various different
16 types of content protection technologies that can be utilized to protect the media
17 content between the physical connector 210 and the output device 212. For
18 example, various content protection types can include Analog Content Protection
19 (ACP), Copy Generation Management System – Analog (CGMS-A), and High-
20 bandwidth Digital Content Protection (HDCP), as will be appreciated by the
21 skilled artisan. In the figure, these various different types of content protection
22 technologies are illustrated at 214.

23 In accordance with the embodiments described herein, a secure link or
24 channel is established between application 202 and one or more downstream
25 components. In the illustrated example, two different secure channels are

1 illustrated, although only one channel of either type need be implemented.
2 Specifically, in one embodiment, a secure communication channel 216 is
3 established between application 202 and driver software 206. In another
4 embodiment, a secure channel 218 is established between application 202 and
5 hardware 208.

6 In the illustrated and described embodiment, channel 216 can be utilized
7 primarily for communication between the application and the driver. Such
8 communication can include communicating *commands*, *status requests* and
9 associated *statuses* between the application and the driver. Channel 218 can be
10 used not only for this type of communication, but for providing the actual media
11 data, such as audio and video data, to the hardware.

12 In the context of this document, a *command* is an instruction and can
13 comprise, among other characteristics, the following characteristics. First, a
14 command can alter or otherwise impact the configuration of the hardware 208.
15 Alternately or additionally, a command can signal the hardware to process the
16 media content in some specific way.

17 A *status request* can effectively retrieve, from the hardware, two sets of
18 information. First, a status request can retrieve information that pertains to any
19 commands that were sent from the application to a downstream component. For
20 example, if the application sent a command to perform some task—for example
21 setting a particular type of content protection technology—an associated status
22 request might retrieve information as to whether the task was performed. Second,
23 a status request can retrieve information that is not necessarily associated with
24 whether a command was sent. For example, a status request might retrieve
25 additional hardware configuration information that might be intrinsic to a

1 particular device, e.g. the type of physical connector that is being utilized. Such
2 status requests can also be used to ascertain which content protection technologies
3 are supported by the hardware for a particular physical connector. For example, a
4 physical connector may be a DVI connector, but there may not be HDCP support
5 in the graphics hardware. As such, the application might then make an intelligent
6 choice to either play a down-sampled or lower resolution version of the video data
7 or perhaps not play the video data at all.

8 9 Establishing a Secure Channel

10 In the illustrated and described embodiment, a one-way trust model and key
11 transport protocol are utilized to establish trust between two endpoints of a
12 communication chain. Using key transport carries with it a couple of advantages,
13 as will be appreciated by the skilled artisan. First, with key transport, the
14 responsibility for creating entropy is left with the entity that establishes trust—
15 here application 202. Second, key transport provides the ability to do what is
16 called “pass through”. Pass through can allow content data to be encrypted all of
17 the way along the chain to the graphics hardware. This, in turn, can allow the
18 graphics hardware to decrypt the content data, decode the content data into an
19 associated surface, and then process the content data onto a display, as will be
20 appreciated by the skilled artisan.

21 Using channel 216 as an example, a one-way authenticated channel is
22 established between application 202 and driver software 206 using, for example,
23 Public Key Infrastructure (PKI) techniques. More specifically, in this example, a
24 channel key in the form of a data integrity key is established between application
25 202 and driver software 206. Once the channel key is established, commands and

1 status messages can be passed unencrypted and message authentication codes or
2 MACs can be created with the data integrity key.

3 As an example, consider the following in connection with Fig. 3. To
4 establish trust with a downstream component, application 202 first calls driver
5 software 206 and receives from the driver software a random number r_{GH} and a
6 digital certificate $Cert_{GH}$ associated with the hardware 208 (Fig. 2). When the
7 application 202 receives the digital certificate $Cert_{GH}$, the application processes the
8 certificate to verify that the signature on the digital certificate is provided by a
9 trusted entity and to retrieve the public key (P_{GH}) of the driver software 206 (in the
10 event channel 216 is utilized) or graphics hardware 208 (in the event channel 218
11 is utilized). The public key P_{GH} can now be used by the application to encrypt
12 data that can only be decrypted by an associated private key which is kept secure
13 by the driver software or the graphics hardware. The process just described
14 establishes trust between the application and the driver software.

15 After trust is established, the application can now set up the key transport.
16 In this example, application 202 concatenates the random number r_{GH} provided by
17 the driver software, a data integrity key k_{DI} (also called a session key), a random
18 starting status sequence number (status_start) and a random starting command
19 sequence number (command_start), encrypts the concatenation of values using the
20 public key P_{GH} of the driver software, and sends the encrypted data to the driver
21 software 206. The process just described establishes the session key that is
22 utilized between the application and the driver software. The random starting
23 status sequence number (status_start) and random starting command sequence
24 number (command_start) are utilized to ensure that if a particular message is lost,
25 whether the message is a status request or a reported status (in the case of the

starting status sequence number) or a command (in the case of the starting command sequence number), the intended message receiver can ascertain this fact. More specifically, each time a command is sent by the application, the random starting command sequence number is incremented by one. Similarly, each time a status message is sent, the sending entity increments the starting status sequence number by one. On the receiving end, the receiver of the message can ascertain whether a message has been lost by simply checking to see whether the sequence numbers occur in order. If there is a missing sequence number, then an associated message has been lost and appropriate action can be taken. In this manner, using random starting sequence numbers ensures the integrity and the order of the messages.

In one example, the values and descriptions set forth in the table just below can be utilized in the process just described.

Value	Description
r_{GH}	Random 32-bit random number generated by the graphics driver.
$Cert_{GH}$	Variable length digital certificate of graphics hardware.
$P_{GH}(r_{GH}, k_{DI}, status_start, command_start)$	Concatenation of 32-bit random number from the graphics driver, 128-bit random data integrity session key, 32-bit random starting status sequence number and 32-bit random starting command sequence number encrypted with the public key of the graphic hardware. This value is 2048 bits long.

At this point in the process, a secure channel has been established between the application and the driver software. The discussion that follows just below

1 describes two portions of the inventive protocol—the command protocol and the
2 status protocol.

3 4 Command Protocol

5 The command portion of the protocol enables the system's hardware, such
6 as hardware 208 (Fig. 2), to be instructed by the application to turn on and in some
7 instances configure a particular content protection technology such as HDCP or
8 CGMS-A. In the illustrated and described embodiment, command messages are
9 contained in an envelope that has two sections: a data section and a message
10 authentication code or MAC section.

11 In the illustrated and described embodiment, the data section of the
12 command message contains a command and an associated command sequence
13 number. The message authentication code or MAC section contains what can be
14 considered as a keyed hash of the command and associated command sequence
15 number. The key that is used to produce the keyed hash is the data integrity key
16 k_{DI} described above.

17 As an example, consider Fig. 4, where application 202 is shown to be
18 sending a command message to the driver software 206 of the form:

19
20 Command, $MAC_{k_{DI}}(\text{Command})$

21
22 where the "Command" portion of the message includes not only an
23 associated command from the application, but the associated command sequence
24 number as well. The " $MAC_{k_{DI}}(\text{Command})$ " portion of the message comprises a
25 MAC of the "Command" portion that was computed by the application. When the

driver software 206 receives the command message, the driver software can run the MAC algorithm over the command portion of the message using the negotiated session key k_{DI} . Once the MAC algorithm has been run over the command portion of the command message, the driver software can compare this value with the MAC value received in the command message. If the two values compare favorably, then the integrity of the command message is established.

In one example, the values and descriptions set forth in the table just below can be utilized in the process just described.

Value	Description
Command	Variable length command data.
$MAC_{k_{DI}}(\text{Command})$	128-bit MAC of the command data and command sequence number using the data integrity session key k_{DI} .

Status Protocol

In the illustrated and described embodiment, status messages are contained in an envelope that has two sections: a data section and a message authentication code or MAC section. In this example, the message sender calculates a MAC of the status data using the session key k_{DI} and an associated MAC algorithm.

As an example of how the status protocol can be implemented, consider Fig. 5. There, application 202 makes a request for status information in the form:

r_{APP} , Status Request

1 where r_{APP} is a random number generated by the application and “Status
2 Request” is the actual request for status information. In response to receiving the
3 status request from the application, driver software 206 ascertain whatever status
4 information is appropriate and provides the status information, along with the
5 appropriate status sequence number, into the data section of the envelope. In
6 addition, the driver software also computes a MAC of the status information along
7 with the random number provided by the application. The status message is then
8 provided back to the application in the following form:

$$r_{APP}, \text{Status}, \text{MAC}_{kDI}(r_{APP}, \text{Status})$$

12 where “Status” includes not only the status information, but the status
13 sequence number as well. The $\text{MAC}_{kDI}(r_{APP}, \text{Status})$ portion of the message
14 comprises a MAC of the random number provided by the application and the
15 status information. When the application receives the status message, it can
16 compute a MAC of the random number and the status information by running the
17 MAC algorithm over the random number and the status information and then can
18 compare this value with the value in the $\text{MAC}_{kDI}(r_{APP}, \text{Status})$ section of the
19 message. By doing this, the application can confirm that the status message is
20 genuine. The application can then use the contents of the message to verify that
21 the status message is in the proper sequence (by virtue of the sequence number)
22 and that the status information pertains to the status request that was previously
23 sent by the application (by virtue of the random number r_{APP}).

24 In one example, the values and descriptions set forth in the table just below
25 can be utilized in the process just described.

Value	Description
r_{APP}	128-bit random number generated by the application.
Status	Variable length status data.
$MAC_{kDI}(r_{APP}, \text{Status})$	128-bit MAC of the status data, status sequence number and the random number provided by the application using the data integrity session key k_{DI} .

Exemplary Method

Fig. 6 illustrates steps in a method in accordance with one embodiment. The method can be implemented in connection with any suitable hardware, software, firmware or combination thereof. Additionally, in some software implementations, aspects of the method can be implemented as computer-readable instructions that reside on one or more computer-readable media and which are executed to perform the described methodology. In this example, the illustrated method is separated into two sections—one of which is labeled “Application”, the other of which is labeled “Downstream Component”. This is done to illustrate which entities perform the acts about to be described. Accordingly, those acts that appear under the label “Application” are performed by the application, and those acts that appear under the label “Downstream Component” are performed by a component that resides downstream of the application. In one embodiment, the downstream component can comprise driver software or a software representation of associated hardware that processes media data. In another embodiment, the downstream component can comprise hardware componentry itself.

1 One exemplary system that can implement the method about to be
2 described is shown and described in connection with Figs. 2-5 above. It is to be
3 appreciated and understood that other systems can be utilized without departing
4 from the spirit and scope of the claimed subject matter.

5 Step 600 ascertains a public key associated with a component downstream
6 from a playback application. Examples of how this can be done are provided
7 above. Step 602 uses the public key to establish a secure channel between the
8 playback application and the downstream component. Collectively, steps 600 and
9 602 establish trust between the playback application and the downstream
10 component. One example of how this trust can be established is described in
11 connection with Fig. 3 above. Other ways can, of course, be used without
12 departing from the spirit and scope of the claimed subject matter.

13 Step 604 uses the secure channel to send a command message to the
14 downstream component. Any suitable command message structure can be utilized
15 with but one example being given above in connection with Fig. 4. In addition,
16 any suitable and/or appropriate types of commands can be sent from the
17 application to the downstream component. In one embodiment, commands that
18 can be sent direct the downstream component to enable a particular type of content
19 protection technology to protect media content or data that is to be processed by
20 the system.

21 Step 606 receives the command message via the secure channel. This step
22 is performed by the downstream component to which the command message is
23 sent. Step 608 then verifies the integrity of the command message. Any suitable
24 method can be used to verify the integrity of the command message with but one
25 example being given in connection with Fig. 4. Once the integrity of the

1 command message is verified, step 610 implements an associated command
2 contained in the command message.

3 Step 612 uses the secure channel to send a status request to the downstream
4 component. Any suitable type of status request structure can be utilized with but
5 one example being given above in connection with Fig. 5. Step 614 receives the
6 status request via the secure channel and step 616 responsively prepares a status
7 message and sends the status message to the application using the secure channel.
8 Any suitable status message structure can be utilized with but one example being
9 given above in connection with Fig. 5.

10 Step 618 receives the status message via the secure channel and step 620
11 verifies the integrity of the status message. Any suitable method can be used to
12 verify the integrity of the status message, with but one example being given above.

13 Exemplary Implementation of Protocol on a Windows® Operating 14 System

15 System Overview

16 Fig. 7 illustrates a system, generally at 700, which can be utilized to
17 implement the secure channel protocol described above. In the discussion that
18 follows, the inventive protocol is referred to as the Certified Output Protection
19 Protocol or "COPP". In this example, Microsoft® DirectShow®, which is an
20 architecture for streaming media, is used as an example. As will be appreciated by
21 the skilled artisan, Microsoft® DirectShow® provides for high quality playback of
22 multimedia streams and supports a wide variety of formats including Advanced
23 Streaming Format (ASF), Motion Picture Experts Group (MPEG), Audio-Video
24 Interleaved (AVI), MPEG Audio Layer-3 (MP3), and WAV files.

1 To achieve throughput that is necessary for streaming video and audio data,
2 DirectShow® uses Microsoft® DirectDraw® and Microsoft® DirectSound® to
3 render data efficiently to the system's sound and graphics cards. DirectDraw® is a
4 display component of DirectX® that allows software designers to directly
5 manipulate display memory, hardware blitters, hardware overlays, and flipping
6 surfaces, as will be appreciated by the skilled artisan. DirectDraw® provides
7 device-independent access to the device-specific display functionality in a direct
8 32-bit path. A 64-bit path also exists in the 64-bit versions of Windows®.
9 DirectDraw® calls important functions in a driver that accesses the display card
10 directly, without the intervention of the Windows® graphics device interface
11 (GDI) or the device-independent bitmap (DIB) engine.

12 In DirectShow®, synchronization is achieved by encapsulating the
13 multimedia data in time-stamped media samples. To handle the variety of
14 different media sources, formats and hardware devices, DirectShow® uses a
15 modular architecture in which operating system components called *filters* can be
16 mixed and matched to provide support for many different scenarios. In the
17 illustrated example, these filters can be assembled into a filter graph 704 that is
18 utilized by a playback application 702 to cause multimedia data to be processed
19 and rendered by a hardware device.

20 In this particular example, filter graph 704 comprises a source filter 706
21 that reads media content and a decoder 708 (e.g. a WMV or MPEG-2 decoder) that
22 decompresses the content. A video rendering component 710 receives the
23 decompressed content or data and understands the hardware on which the content
24 is to be rendered. In this particular example, there are two different types of
25 rendering components that can be utilized—a VMR7 (Video Mixing Renderer

1 Version 7 written to the DirectX version 7 SDK) and a VMR9 (Video Mixing
2 Renderer Version 9 written to the DirectX version 9 SDK) component.

3 In the illustrated example, a user mode DirectDraw® component 712 is
4 provided and is a system-supplied dynamic-link library (DLL) that is loaded and
5 called by DirectDraw® applications. This component provides hardware
6 emulation, manages the various DirectDraw® objects, and provides display
7 memory and display management services, as will be appreciated by the skilled
8 artisan.

9 In addition, a kernel mode DirectDraw® component 714 is provided and is
10 an integral part of win32k.sys. Kernel mode component 714 is the system-
11 supplied graphics engine that is loaded by a kernel-mode display driver. This
12 component performs parameter validation for the driver, thus making it easier to
13 implement more robust drivers. A DirectDraw® graphics driver 716 is provided
14 and is typically implemented by a third party hardware vendor.

15 In operation, the protected content typically has some set of policies
16 associated with it. The policies can establish the type and level of content
17 protection that is to be used when the associated content is rendered. For example,
18 if the content is going to be played out over a DVI connector, then a policy might
19 establish that HDCP has to be enabled or the content will not be played.

20 In the present example, application 702 is aware of the policies that are
21 associated with the content. In practice, the source filter 706 can include a
22 component that is responsible for ascertaining the policy associated with the
23 content and then notifying the application.

User Mode Application Program Interface (API)

Once the policies have been extracted and the implications of the policies ascertained, system 700 first ascertains whether the policy can be enforced in this specific system. The way that this is accomplished, in this particular example, is through a user mode application program interface (API) called IAMCertifiedOutputProtection which is exposed by the video rendering component 710. In accordance with one embodiment, the IAMCertifiedOutputProtection interface consists of the following methods which can be called by the application:

```
typedef struct _AMCOPPSignature {  
    BYTE    Signature[256];  
} AMCOPPSignature;  
  
typedef struct _AMCOPPCCommand {  
    GUID     mackDI;           // 16 bytes  
    GUID     guidCommandID;    // 16 bytes  
    DWORD    dwSequence;      // 4 bytes  
    DWORD    cbSizeData;       // 4 bytes  
    BYTE     CommandData[4056]; // 4056 bytes (4056+4+4+16+16 = 4096)  
} AMCOPPCCommand, *LPAMCOPPCCommand;  
  
typedef struct _AMCOPPStatusInput {  
    GUID     rApp;             // 16 bytes  
    GUID     guidStatusRequestID; // 16 bytes  
    DWORD    dwSequence;      // 4 bytes  
    DWORD    cbSizeData;       // 4 bytes  
    BYTE     StatusData[4056];  // 4056 bytes (4056+4+4+16+16 = 4096)  
} AMCOPPStatusInput, *LPAMCOPPStatusInput;  
  
typedef struct _AMCOPPStatusOutput {  
    GUID     mackDI;           // 16 bytes  
    DWORD    cbSizeData;       // 4 bytes
```

```

1      BYTE    COPPStatus[4076];    // 4076 bytes (4076+16+4 = 4096)
2  } AMCOPPStatusOutput, *LPAMCOPPStatusOutput;
3
4  IAMCertifiedOutputProtection
5
6  HRESULT KeyExchange (
7      [out]    DWORD* pdwRandom          32-bit random number
8                                          generated by Graphics Driver
9                                          Graphics Hardware
10                                         certificate, memory released
11                                         by CoTaskMemFree
12                                         [out]    DWORD* pdwLengthCertGH; Length of Graphics Hardware
13                                         certificate
14
15  HRESULT SessionSequenceStart(
16      [in]      AMCOPPSignature          Concatenation of 32-bit random
17      *pSig);                             number from the graphics
18                                         driver, 128-bit random data
19                                         security session key, 128-bit
20                                         random data integrity session
21                                         key, 32-bit random starting
22                                         status sequence number and 32-
23                                         bit random starting command
24                                         sequence number encrypted with
25                                         the public key of the graphic
                                         hardware. This value is 2048
                                         bits long.
26
27  HRESULT ProtectionCommand(
28      [in]      const AMCOPPCmd * cmd,    Encrypted command
29
30  HRESULT ProtectionStatus(
31      [in]      const
32      AMCOPPStatusInput*          Description of feature we are
33      pInput,                     requesting status on.
34      [out]     AMCOPPStatusOutput*    Returned status information
35      pOutput);

```

In the above example, the AMCOPPCmd structure is the data structure that the application uses to send a command to a downstream component and would be used in an exchange such as the one illustrated in Fig. 4. The AMCOPPStatusInput structure is the data structure that the application uses to make a status request and would be used in an exchange such as the one illustrated on the left side of Fig. 5. The AMCOPPStatusOutput structure is the data structure

1 that the graphics driver uses to provide status information back to the application
2 and would be used in an exchange such as the one illustrated on the right side of
3 Fig. 5.

4 5 Kernel Mode DDI (DX-VA)

6 The DX-VA Certified Output Protection Protocol (COPP) provides a means
7 for passing commands from User mode applications, such as playback application
8 702, to a driver such as graphics driver 716, and for returning status from the
9 driver to the user mode applications. As noted above, the channel through which
10 the commands and status flow is protected.

11 In this particular implementation, the Content Output Protection Protocol
12 utilizes a DX-VA COPP device which is described in more detail below. In the
13 discussion just below, the following terminology is used:

- 14
15 • **Video Stream:** A Video Stream is the principal image data that
16 comprises a displayed video frame. Pixels from the Video Stream
are always opaque and do not contain any per-pixel alpha data.
- 17
18 • **Video Sub-stream:** A Video Sub-stream is a channel of auxiliary
19 image data that is required to be composited with a sample from the
20 Video Stream prior to display of the combined Video frame. Examples of Video Sub-streams are NTSC Closed Caption images, DVD Sub-Picture images and PAL Teletext images. Another characteristic of Video Sub-streams is that they contain limited color range and per-pixel alpha information.
- 21
22 • **Video Session:** A video session consists of a Video Stream and
23 possibly one or more Video Sub-streams which get combined with
24 the video stream before it is displayed on an output device connected
25 to the computer. There may be several Video Sessions active on the
computer at any particular time, there may also be several Video
Sessions active within a single process.

- 1 • **DX-VA COPP Device:** A new DX-VA device that supports
2 receiving commands and status requests via the Content Output
3 Protection Protocol.
- 4 • **Connector:** The physical output connection between the graphics
5 hardware and a display device.
- 6 • **Protection Type:** The type of protection that can be applied to the
7 signal being passed through a connector; more than one type of
8 protection can be applied to a single connector.
- 9 • **Protection Level:** The level of protection that is applied to signal
10 being passed through a connector. This value is dependent on the
11 protection type – some protections types, for example HDCP, have
12 only two protection levels, these being “on” and “off”.

13 Design Overview

14 In accordance with this embodiment, a single component known as a DX-
15 VA COPP device is created for each video session. In the Fig. 7 example, two
16 exemplary DX-VA COPP devices are shown within driver 716. The DX-VA
17 COPP device represents an end point for COPP commands and status requests. In
18 addition, a DX-VA COPP device can virtualize the protection settings of a physical
19 connector. It is possible that a single physical connector can support multiple
20 content protection types. For example, an S-Video connector may support Analog
21 Content Protection and well as CGMS-A Line20 protection and CGMS-A Line21
22 protection. The COPP sample device driver code that appears later in this
23 document shows one example of how a COPP device can be implemented by a
24 driver.

25 In the illustrated and described embodiment, each DX-VA COPP device is
configured to act appropriately when multiple video sessions are active on the
computer. In this case, the driver 716 is asked to create multiple instances of the

1 COPP DX-VA device as different processes try to configure the output display
2 adapter settings via COPP. In the present example, two instances of a DX-VA
3 COPP devices have been created by driver 716.

4 5 Local and Global Reference Counts

6 For each physical connector that supports content protection, driver 716
7 should maintain a global reference count for each type and for each level. In the
8 illustrated and described embodiment, driver 716 comprises a global state
9 component that maintains this global state.

10 When a COPP DX-VA device is created for a video session, the device
11 should contain a local reference count for each protection type at each protection
12 level. This is represented by the current protection level component inside each of
13 the devices. A default level counter for the protection type should be set to the
14 value 1, and all the other protection level counters for that type should be set to the
15 value zero.

16 When a video session sets a new protection level for a particular protection
17 type, the reference count for the current protection level is decremented and the
18 reference count for the new protection level is incremented. Corresponding
19 changes are also made to the global reference level counters.

20 Whenever any of the global level counters change, the driver should inspect
21 all the counters for that connector and ensure that the protection level is set to
22 level that corresponds with highest level counter whose value is greater than zero.
23 In the present example, there are two different protection levels for an associated
24 physical connector—level A and level B. Notice that the global state component
25 indicates that level B protection is applied to the physical connector. As such,

1 level B is the higher of the protection levels that can be applied to the physical
2 connector. If DX-VA COPP device 2 terminates, then the global state component
3 will apply protection level A to the physical connector. Exemplary code samples
4 below illustrate how this might be done.

5 While the global reference counter is greater than zero, content protection
6 should be applied to the output connector. As soon as the global reference counter
7 reaches zero, content protection should be removed from the output connector.
8 Whenever a driver receives a DestroyDevice call, the global reference counter
9 should be decremented by the current level of the devices local reference counter.
10 The connector protection should only be removed if the global reference counter
11 for that connector reaches zero.

12 It should be noted that DestroyDevice may be called while the device still
13 has local reference counter greater than zero, one example of this would be if the
14 user mode process terminated abnormally.

16 Video Session Initialization

17 In the illustrated and described embodiment, video sessions are tied to a
18 particular physical output connector. A video session, at the display driver level, is
19 initialized by the following steps:

- 21 • **GetMoCompGuids** is called to determine if a driver supports a DX-
VA COPP device for the output connector.
- 22 • Assuming the driver does support a DX-VA COPP device for the
output connector **CreateMoComp** is called to create the DX-VA
23 COPP device.
- 24 • **RenderMoComp** is called with a
DXVA_COPPGetCertificateLengthFnCode function code to
25

determine the length of the variable length graphics hardware certificate that should be used for this session.

- **RenderMoComp** is called with a DXVA_COPPKeyExchangeFnCode function code to pass the Video Sessions Application certificate to the driver. The driver returns a 128 bit privacy key and a copy of its variable length graphics hardware certificate.
- Finally, **RenderMoComp** is called with a DXVA_COPPSequenceStartFnCode function code to indicate that the Video Session has started with a 32 bit Command Sequence start code and 32 bit Status Sequence start code.

Details for the individual DX-VA RenderMoComp function codes are given below.

Handling Device Lost Conditions

The following conditions will cause **DestroyMoComp** entry point in the driver to be called while output content protection is still enabled for the specific Video Session:

- Display Mode changes;
- Attaching or detaching a monitor from the Windows® desktop;
- Entering a full screen DOS box;
- Starting any DDraw or D3D exclusive mode application;
- Fast User Switching;
- Adding or removing monitors from the Windows® desktop;
- Locking the workstation or pressing Alt-Ctrl-Del on Windows® XP professional edition;
- Attaching to the workstation via “Remote Desktop Connection”;
- Entering a power saving mode, for example “Suspend” or “Hibernate”; and
- Unexpected application termination, for example a “page fault”.

When this occurs the driver should decrement the global protection level count by the current local protection level count for the device being destroyed.

1 The changed global protection level should be examined and the protection
2 applied to the output connector should be adjusted accordingly.

4 DX-VA Commands

5 The following discussion describes a number of commands or functions
6 that are supported by the graphics driver and, more particularly, the individual
7 DX-VA COPP devices.

9 **COPPOpenVideoSession**

10 This function initializes the COPP DX-VA device that is being used for a
11 particular video session.

```
12 HRESULT DXVA_COPPDeviceClass::COPPOpenVideoSession(  
13     DWORD dwDeviceID  
14 );
```

15 **COPPCloseVideoSession**

16 This function terminates the COPP DX-VA device that is being used for the
17 video session.

```
18  
19 HRESULT DXVA_COPPDeviceClass::COPPCloseVideoSession();  
20
```

21 It is possible that this function may be called while there is output
22 protection still applied by the video session. The driver should undo the protection
23 settings of this DX-VA COPP device and adjust the global protection settings
24 accordingly.
25

COPPGetCertificateLength

This is the first function call that a newly created COPP DX-VA device will receive. It is used by the video rendering component 710 to query the length, in bytes, of the certificate used by the graphics hardware. The driver should return the correct length in the *lpOutputData* field of the DD_RENDERMOCOMPPDATA structure.

```
HRESULT DXVA_COPPDeviceClass::COPPGetCertificateLength(  
    [out] DWORD* pCertificateLength  
);
```

COPPKeyExchange

This is the second function call that a new created COPP DX-VA device will receive. The driver should return the variable length digital certificate of the graphics hardware. The *lpOutputData* field of the DD_RENDERMOCOMPPDATA structure points to a DXVA_COPPKeyExchangeOutput structure; the DXVA_COPPVariableLengthData field will have the required amount of space for the certificate as indicated in the previous COPPGetCertificateLength DXVA command.

```
typedef struct _DXVA_COPPSignature {  
    BYTE    Signature[256];  
} DXVA_COPPSignature, *LPDXVA_COPPSignature;  
  
typedef struct _DXVA_COPPVariableLengthData {  
    DWORD    Size;  
    BYTE    Data[4];  
} DXVA_COPPVariableLengthData, *LPDXVA_COPPVariableLengthData;  
  
typedef struct _DXVA_COPPKeyExchangeOutput {  
    DWORD    RandomNumber;  
    DXVA_COPPVariableLengthData Certificate;  
} DXVA_COPPKeyExchangeOutput, *LPDXVA_COPPKeyExchangeOutput;
```

```

1  HRESULT DXVA_COPPDeviceClass::COPPKeyExchange(
    [out] DWORD* pdwRandom,
2  [out] DXVA_COPPVariableLengthData* pGHCertificate
    );
3

```

4 **COPPSequenceStart**

5 This will be the third functional call that the DX-VA COPP device will
6 receive. The driver will receive a concatenation of a 128-bit random data integrity
7 session key, 32-bit random starting status sequence number, and 32-bit random
8 starting command sequence number encrypted with the public key of the graphic
9 hardware. The driver should return E_UNEXPECTED if it receives this function
10 call before it receives the COPPKeyExchange function call or if it receives a
11 subsequent COPPSequenceStart call.

```

12
13 HRESULT DXVA_COPPDeviceClass::COPPSequenceStart(
    [in] DXVACOPPSignature* pSeqStartInfo
14 );
15

```

15 **COPPCommand**

16 A COPP command can comprise an instruction from the video session to
17 change a protection level on the physical connector associated with the DX-VA
18 COPP device. The driver should be able to support multiple video sessions all
19 playing back content through the same physical connector, as noted above.

```

20
21 typedef struct _DXVACOPPCommand {
    GUID      macKDI;           // 16 bytes
    GUID      guidCommandID;    // 16 bytes
22     DWORD    dwSequence;      // 4 bytes
    DWORD     cbSizeData;       // 4 bytes
23     BYTE     CommandData[4056]; // 4056 bytes (4056+4+4+16+16 = 4096)
} DXVA_COPPCommand, *LPDXVA_COPPCommand;

```

```

24 typedef struct _DXVA_COPPSetProtectionLevelCmdData {
    DWORD     ProtType;
25     DWORD     ProtLevel;

```

```

} DXVA_COPPSetProtectionLevelCmdData;
1
DEFINE_GUID(DXVA_COPPSetProtectionLevel,
2
    0x9bb9327c, 0x4eb5, 0x4727, 0x9f, 0x00, 0xb4, 0x2b, 0x09, 0x19, 0xc0, 0xda);

HRESULT DXVA_COPPDeviceClass::COPPCommand(
3
    [in] DXVA_COPPCommand* pCommand
4
    );

```

The driver should return ***E_UNEXPECTED*** if it receives this function call before it receives the COPPKeyExchange or COPPSequenceStart function calls. Also, the driver should ensure that the parameters passed to it are valid for the given physical connector being used. The driver should return ***E_INVALIDARG*** if one or more of the parameters passed in the COPPCommand are not valid.

The table below lists the set of DX-VA COPP commands that a driver can expect to see in accordance with this embodiment.

GUID	Interpretation of CommandData Parameter	Description
DXVA_COPPSetProtectionLevel 9bb9327c-4eb5-4727-9f00b42b0919c0da	DXVA_COPPSetProtectionLevelCmdData Valid values for ProtType member variable are: COPP_ProtectionType_ACP COPP_ProtectionType_CGMSA COPP_ProtectionType_HDCP	Sets the specified protection type to the specified protection level.

Defined COPP Commands

COPPQueryStatus

A COPPQueryStatus is a request from the video session to retrieve information about the physical connector being used, the type of protection that can be applied to the content being transmitted through the physical connector, and the current protection level that is active on the physical connector. The same restriction about the order of the calls applies here: the driver should return

E_UNEXPECTED if it receives this call before either COPPKeyExchange or COPPSequenceStart calls for this VideoSession.

```
typedef struct _DXVA_COPPStatusInput {
    GUID      rApp;                // 16 bytes
    GUID      guidStatusRequestID; // 16 bytes
    DWORD     dwSequence;          // 4 bytes
    DWORD     cbSizeData;          // 4 bytes
    BYTE      StatusData[4056];    // 4056 bytes (4056+4+4+16+16 = 4096)
} DXVA_COPPStatusInput, *LPDXVA_COPPStatusInput;

typedef struct _DXVA_COPPStatusOutput {
    GUID      macKDI;              // 16 bytes
    DWORD     cbSizeData;          // 4 bytes
    BYTE      COPPStatus[4076];    // 4076 bytes (4076+16+4 = 4096)
} DXVA_COPPStatusOutput, *LPDXVA_COPPStatusOutput;

HRESULT DXVA_COPPDeviceClass::COPPQueryStatus(
    [out] DXVA_COPPStatusInput* pInput,
    [out] DXVA_COPPStatusOutput* pOutput
);
```

The table below lists a set of DX-VA status requests that a driver should expect to see in accordance with this embodiment.

GUID	Input Data	Returned Data	Description
DXVA_COPPQuery ConnectorType 81d0bfd5-6afe-48c2- 99c0-95a08f97c5da	None	ΓAPP (GUID 16 bytes) plus DWORD (4 bytes) COPP_ConnectorType _Unknown COPP_ConnectorType _VGA COPP_ConnectorType _SVideo COPP_ConnectorType _CompositeVideo COPP_ConnectorType _ComponentVideo COPP_ConnectorType _DVI COPP_ConnectorType _HDMI COPP_ConnectorType _LVDS	Returns the a value that identifies the type of physical connector the Video Session is currently using.
DXVA_COPPQuery ProtectionType 38f2a801-9a6c- 48bb-9107- b6696e6f1797	None	ΓAPP (GUID 16 bytes) plus DWORD (4 bytes) COPP_ProtectionType _Unknown COPP_ProtectionType _None COPP_ProtectionType _ACP COPP_ProtectionType _CGMSA COPP_ProtectionType _HDCP	Returns the available protection mechanism for the physical connector. These values can be "or'ed" together if multiple protection types are supported on a physical connector.
DXVA_COPPQuery LocalProtectionLevel b2075857-3eda- 4d5d-88db- 748f8c1a0549	DWORD (4 Bytes) – Identifies the protection type that we require status upon. COPP_ProtectionT ype_ACP COPP_ProtectionT ype_CGMSA COPP_ProtectionT ype_HDCP	ΓAPP (GUID 16 bytes) plus 2x DWORD (8 bytes) The first DWORD identifies the protection level of the specified protection type, the actual values returned depend on the protection type for this connector. The second DWORD identifies the integrity of the connection between the computer and the display device. The value returned could be either:	Returns the currently set protection level for the Video Session.

		COPP_LinkActive COPP_LinkLost	
DXVA_COPPQuery GlobalProtectionLevel 1957210a-7766-452a-b99a-d27aed54f03a	DWORD (4 Bytes) – Identifies the protection type that we require status upon. COPP_ProtectionType_ACP COPP_ProtectionType_CGMSA COPP_ProtectionType_HDCP	ΓAPP (GUID 16 bytes) plus 2x DWORD (8 bytes) The first DWORD identifies the protection level of the specified protection type, the actual values returned depend on the protection type for this connector. The second DWORD identifies the integrity of the connection between the computer and the display device. The value returned could be either: COPP_LinkActive COPP_LinkLost	Returns the currently set protection level for the physical connector.

Defined COPP Status Requests

DDI Mapping for User Mode APIs

In this particular implementation example, and for compatibility with the DDI infrastructure for Windows® operating systems, the API described earlier in this document must be “mapped” to the existing DDI for DirectDraw and DirectX VA. These mapping operations help to avoid having to make significant changes to kernel mode components, as will be appreciated by the skilled artisan. Thus, this section and the discussion below describes the COPP interface mapping to the existing DirectDraw and DX-VA DDI.

The DX-VA DDI is itself split into two functional groups: the “DX-VA container” and the “DX-VA device.” The purpose of the DX-VA container DDI group is to determine the number and capabilities of the various DX-VA devices contained by the display hardware. Therefore, in this implementation example, a DX-VA driver can only have a single container, but it can support multiple DX-VA devices.

1 The rest of this section describes how the interface described above is
2 mapped to the DX-VA device DDI. It should be noted that unlike other DX-VA
3 DDI's, COPP devices never reference any video memory surfaces.

5 Calling the DDI from a User-Mode Component

6 The sequence of steps to use the DDI from a user-mode component, such as
7 the video rendering component 710 (Fig. 7) is as follows:

- 8 1. Call ***GetMoCompGuids*** to get the list of DX-VA devices supported
9 by the driver.
- 10 2. ***CreateMocomp*** is called to establish a COPP device for the video
11 session, where the COPP device GUID is defined as follows:

```
12           DEFINE_GUID(DXVA_COPPDevice,  
13           0xd2457add,0x8999,0x45ed,0x8a,0x8a,0xd1,0xaa,0x04,0x7b,0xa4,0xd5);
```

- 14 3. The video rendering component then calls the COPP device's
15 ***RenderMocomp*** with varying function parameter values depending
16 on the type of COPP operation being requested.
- 17 4. When the video rendering component no longer needs to perform
18 any more COPP operations, it calls ***DestroyMocomp***.
- 19 5. The driver releases any resources used by the COPP device.

20 COPPOpenVideoSession

21 This method maps directly to a ***CreateMoComp*** method of the
22 ***DD_MOTIONCOMPCALLBACKS*** structure, where the GUID is the COPP
23 Device GUID, *pUncompData* points to a structure that contains no data (all zeros),
24 and *pData* points to NULL.

25 It should be appreciated and understood that if a driver supports accelerated
decoding of compressed video, the video rendering component will call the driver

to create two DX-VA devices—one to perform the actual video decoding work as defined by the DirectX VA Video Decoding specification, and the other to be used strictly for applying protection to output connectors following the Content Output Protection Protocol.

Example: Mapping CreateMoComp to COPPOpenVideoSession

The sample code provide just below shows how a driver should map the *CreateMoComp* DDI call into calls to *COPPOpenVideoSession*. The sample code shows only how the *CreateMoComp* function is used for COPP. If a driver supports other DX-VA functions, such as decoding MPEG-2 video streams, the sample code below can be extended to include processing of additional DX-VA GUIDs.

```
DWORD APIENTRY
CreateMoComp(
    LPDDHAL_CREATEMOCOMPDATA lpData
)
{
    AMTRACE((TEXT("CreateMoComp ")));

    if (!ValidDXVAGuid(lpData->lpGuid))
    {
        DbgLog((LOG_ERROR, 1,
            TEXT("No formats supported for this GUID")));
        lpData->ddrVal = E_INVALIDARG;
        return DDHAL_DRIVER_HANDLED;
    }

    if (*lpData->lpGuid == DXVA_COPPDevice)
    {
        DXVA_COPPDeviceClass* lpDev =
            new DXVA_COPPDeviceClass(*lpData->lpGuid, DXVA_DeviceCOPP);

        if (lpDev)
        {
            DWORD DevID = (DWORD)-1;
            lpData->ddrVal = E_INVALIDARG;

            if (DevID != (DWORD)-1)
            {
                lpData->ddrVal = lpDev->COPPOpenVideoSession(DevID);
            }

            if (lpData->ddrVal != DD_OK)
```

```

1         {
2             delete lpDev;
3             lpDev = NULL;
4         }
5     else
6     {
7         lpData->ddRVal = E_OUTOFMEMORY;
8     }
9
10    lpData->lpMoComp->lpDriverReserved1 =
11        (LPVOID) (DXVA_DeviceBaseClass*)lpDev;
12    return DDHAL_DRIVER_HANDLED;
13 }
14
15 lpData->ddRVal = DDERR_CURRENTLYNOTAVAIL;
16 return DDHAL_DRIVER_HANDLED;
17 }

```

Example: Implementing GetMoCompGuids

In addition to the *CreateMoComp* DDI function, the driver should also implement the *GetMoCompGuids* method of the **DD_MOTIONCOMPCALLBACKS** structure. The following sample code shows one possible way of implementing this function in your driver.

```

15 DWORD g_dwDXVNumSupportedGUIDs = 1;
16 const GUID* g_DXVASupportedGUIDs[1] = {
17     &DXVA_COPPDevice
18 };
19
20 DWORD APIENTRY
21 GetMoCompGuids(
22     PDD_GETMOCOMPGUIDSDATA lpData
23 )
24 {
25     DWORD dwNumToCopy;
26
27     if (lpData->lpGuids) {
28         dwNumToCopy = min(g_dwDXVNumSupportedGUIDs,
29                          lpData->dwNumGuids);
30         for (DWORD i = 0; i < dwNumToCopy; i++) {
31             lpData->lpGuids[i] = *g_DXVASupportedGUIDs[i];
32         }
33     }
34     else {
35         dwNumToCopy = g_dwDXVNumSupportedGUIDs;
36     }
37
38     lpData->dwNumGuids = dwNumToCopy;
39     lpData->ddRVal = DD_OK;
40     return DDHAL_DRIVER_HANDLED;

```

COPPGetCertificateLength

This method maps directly to a *RenderMoComp* method of the **D_MOTIONCOMPCALLBACKS** structure, where:

- **dwNumBuffers** is zero.
- **lpBufferInfo** is NULL.
- **dwFunction** is defined as **DXVA_COPPGetCertificateLengthFnCode**.
- **lpInputData** is NULL.
- **lpOutputData** will point to a single DWORD.

Note that for the DX-VA device used for COPP, *RenderMoComp* will be called without calling *BeginMoCompFrame* or *EndMoCompFrame*.

Example: Mapping RenderMoComp to COPPGetCertificateLength

The sample code below shows how a driver should map the *RenderMoComp* DDI call into calls to *COPPGetCertificateLength*. The sample code only shows how the *RenderMoComp* function is used for COPP control. If a driver supports other DX-VA functions, such as decoding MPEG-2 video streams, the sample code can be extended to include processing of additional DX-VA GUIDs.

```
DWORD APIENTRY
RenderMoComp(
    LPDDHAL_RENDERMOCOMPDATA lpData
)
{
    AMTRACE((TEXT("RenderMoComp")));
    DXVA_DeviceBaseClass* pDXVABase =
        (DXVA_DeviceBaseClass*)lpData->lpMoComp->lpDriverReserved1;

    if (pDXVABase == NULL) {
        lpData->ddrVal = E_POINTER;
        return DDHAL_DRIVER_HANDLED;
    }
}
```

```

switch (pDXVABase->m_DeviceType) {
1
case DXVA_DeviceCOPP:
2
    {
        DXVA_COPPDeviceClass* pDXVACopp =
3
            (DXVA_COPPDeviceClass*)lpData->lpMoComp->lpDriverReserved1;

        switch (lpData->dwFunction)
4
        {
            case DXVA_COPPGetCertificateLengthFnCode:
5
                {
                    if (lpData->dwOutputDataSize < sizeof(DWORD))
6
                    {
                        lpData->ddRVal = E_INVALIDARG;
7
                    }
                    else
8
                    {
                        lpData->ddRVal = pDXVACopp->COPPGetCertificateLength(
                            (DWORD*)lpData->lpOutputData);
9
                    }
                }
            break;
10
        }
    }
}

```

COPPKeyExchange

This method maps directly to a *RenderMoComp* method of the **D_MOTIONCOMPCALLBACKS** structure, where:

- **dwNumBuffers** is zero.
- **lpBufferInfo** is NULL.
- **dwFunction** is defined as **DXVA_COPPKeyExchangeFnCode**.
- **lpInputData** is NULL.
- **lpOutputData** will point to a **DXVA_COPPKeyExchangeOutput** data structure.

Note that for the DX-VA device used for COPP, *RenderMoComp* will be called without calling *BeginMoCompFrame* or *EndMoCompFrame*.

Example: Mapping RenderMoComp to COPPKeyExchange

The sample code below shows how a driver should map the *RenderMoComp* DDI call into calls to *COPPKeyExchange*. The sample code

only shows how the *RenderMoComp* function is used for COPP control. If a driver supports other DX-VA functions, such as decoding MPEG-2 video streams, the sample code can be extended to include processing of additional DX-VA GUIDs.

```
case DXVA_COPPKeyExchangeFnCode:
{
    DXVA_COPPKeyExchangeOutput* lpout =
        (DXVA_COPPKeyExchangeOutput*)lpData->lpOutputData;

    lpData->ddRVal = pDXVACopp->COPPKeyExchange(
        &lpout->RandomNumber,
        &lpout->Certificate);
}
break;
```

COPPSequenceStart

This method maps directly to a *RenderMoComp* method of the **D_MOTIONCOMPCALLBACKS** structure, where:

- **dwNumBuffers** is zero.
- **lpBufferInfo** is NULL.
- **dwFunction** is defined as **DXVA_COPPSequenceStartFnCode**.
- **lpInputData** will point to a DXVA_COPPSignature data structure.
- **lpOutputData** is NULL.

Note that for the DX-VA device used for COPP, *RenderMoComp* will be called without calling *BeginMoCompFrame* or *EndMoCompFrame*.

Example: Mapping RenderMoComp to COPPSequenceStart

The sample code below shows how a driver should map the *RenderMoComp* DDI call into calls to *COPPSequenceStart*. The sample code only shows how the *RenderMoComp* function is used for COPP control. If a

1 driver supports other DX-VA functions, such as decoding MPEG-2 video streams,
2 the sample code can be extended to include processing of additional DX-VA
3 GUIDs.

```
4  
5     case DXVA_COPPSignatureStartFnCode:  
6     {  
7         DXVA_COPPSignature* lpin =  
8             (DXVA_COPPSignature*)lpData->lpInputData;  
9         lpData->ddrVal = pDXVACopp->COPPSequenceStart(lpin);  
10    }  
11    break;
```

10 COPPCommand

11 This method maps directly to a *RenderMoComp* method of the
12 **D_MOTIONCOMPCALLBACKS** structure, where:

- 13 • **dwNumBuffers** is zero.
- 14 • **lpBufferInfo** is NULL.
- 15 • **dwFunction** is defined as **DXVA_COPPCommandFnCode**.
- 16 • **lpInputData** will point to a **DXVA_COPPCommand** data structure.
- 17 • **lpOutputData** is NULL.

18 Note that for the DX-VA device used for COPP, *RenderMoComp* will be
19 called without calling *BeginMoCompFrame* or *EndMoCompFrame*.

20 Example: Mapping RenderMoComp to COPPCommand

21 The sample code below shows how a driver should map the
22 *RenderMoComp* DDI call into calls to *COPPCommand*. The sample code only
23 shows how the *RenderMoComp* function is used for COPP control. If a driver
24 supports other DX-VA functions, such as decoding MPEG-2 video streams, the
25 sample code can be extended to include processing of additional DX-VA GUIDs.


```

1
2
3         case DXVA_COPPCCommandFnCode:
4             {
5                 DXVA_COPPCCommand* lpin =
6                     (DXVA_COPPCCommand*)lpData->lpInputData;
7                 lpData->ddRVal = pDXVACopp->COPPCCommand(lpin);
8             }
9             break;
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25

```

COPPQueryStatus

This method maps directly to a *RenderMoComp* method of the **D_MOTIONCOMPCALLBACKS** structure, where:

- **dwNumBuffers** is zero.
- **lpBufferInfo** is NULL.
- **dwFunction** is defined as DXVA_COPPQueryStatusFnCode.
- **lpInputData** will point to a DXVA_COPPStatusInput data structure.
- **lpOutputData** will point to a DXVA_COPPStatusOutput data structure.

Note that for the DX-VA device used for COPP, *RenderMoComp* will be called without calling *BeginMoCompFrame* or *EndMoCompFrame*.

Example: Mapping RenderMoComp to COPPQueryStatus

The sample code below shows how your driver should map the *RenderMoComp* DDI call into calls to *COPPQueryStatus*. The sample code only shows how the *RenderMoComp* function is used for COPP control. If a driver supports other DX-VA functions, such as decoding MPEG-2 video streams, the sample code can be extended to include processing of additional DX-VA GUIDs.

```

24
25
26         case DXVA_COPPQueryStatusFnCode:
27             {
28
29
30
31
32
33
34
35
36
37
38
39
40
41
42
43
44
45
46
47
48
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```

```

1         DXVA_COPPStatusInput* lpin = (DXVA_COPPStatusInput*)
           lpData->lpInputData;
2         DXVA_COPPStatusOutput* lpout = (DXVA_COPPStatusOutput*)
           lpData->lpOutputData;
           lpData->ddRVal = pDXVACopp->COPPQueryStatus(lpin, lpout);
3     }
    break;
4
    default:
        lpData->ddRVal = E_INVALIDARG;
        break;
5    }
    break;
6    }
    return DDHAL_DRIVER_HANDLED;
7 }

```

COPPCloseVideoSession

This method maps directly to a *DestroyMoComp* method of the **DD_MOTIONCOMPCALLBACKS** structure.

Example: Mapping DestroyMoComp to COPPCloseVideoSession

The following sample code shows how a driver should map the *DestroyMoComp* DDI call into calls to *COPPCloseVideoSession*. The sample code shows only how the *DestroyMoComp* function is used for the COPP device. If a driver supports other DX-VA functions, such as decoding MPEG-2 video streams, the sample code can be extended below to include processing of additional DX-VA GUIDs.

```

20  DWORD APIENTRY
    DestroyMoComp(
21      LPDDHAL_DESTROYMOCOMPDATA lpData
    )
    {
22      AMTRACE((TEXT("DestroyMoComp")));
      DXVA_DeviceBaseClass* pDXVABase =
23          (DXVA_DeviceBaseClass*) lpData->lpMoComp->lpDriverReserved1;

      if (pDXVABase == NULL) {
24          lpData->ddRVal = E_POINTER;
          return DDHAL_DRIVER_HANDLED;
25      }
    }

```

```

switch (pDXVABase->m_DeviceType) {
1  case DXVA_DeviceContainer:
   case DXVA_DeviceDecoder:
   case DXVA_DeviceDeinterlacer:
2  case DXVA_DeviceProcAmpControl:
    lpData->ddRVal = DDERR_CURRENTLYNOTAVAIL;
    break;
3
   case DXVA_DeviceCOPP:
4       {
        DXVA_COPPDeviceClass* pDXVADev = (DXVA_COPPDeviceClass*)pDXVABase;
        lpData->ddRVal = pDXVADev->COPPCloseVideoSession();
5        delete pDXVADev;
        }
6    break;
7 }

```

Conclusion

Various embodiments described above enable media content protection by establishing a secure communication channel and, in some embodiments, a secure data channel, between a device such as a computer running a protected content playback application, and an associated driver, such as a graphics driver, of an associated display device such as a monitor, flat panel LCD, television and the like. Various embodiments address the needs of content providers and application vendors to signal that content protection should be applied to media output by utilizing an output adapter which physically links the computer and its display device as the gateway for the protected video path.

The various embodiments provide a means for securely passing commands from a user mode playback application to a driver, such as a graphics driver, and securely returning status from the driver to the user mode application. The described embodiments establish a secure communication channel and allow a user mode application to instruct an associated driver to enable content output protection on the physical connector between the computer and the display device.

Although the invention has been described in language specific to structural features and/or methodological steps, it is to be understood that the invention

1 defined in the appended claims is not necessarily limited to the specific features or
2 steps described. Rather, the specific features and steps are disclosed as preferred
3 forms of implementing the claimed invention.
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